Programming for mathematical research

Ken Jacobs, Luca Schaffler University of Georgia

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To build intuition;

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To actually prove a theorem (working case-by-case);

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- To build intuition;
- To actually prove a theorem (working case-by-case);
- To generate supporting examples for your theory.

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Good kinds of problems:

- Something that involves enumeration of finitely many cases;
- Visualization;
- Something you know there is an efficient package for already.

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- Visualization;
- Something you know there is an efficient package for already.

Bad kinds of problems:

- something which is too simple;
- complicated expressions with many indeterminates;
- too much work to tell the computer how to do it.

Does the homogeneous polynomial

$$F(X, Y) = 3X^{3}Y + 3X^{2}Y^{2} + X^{4} - 6Y^{4}$$

factor over $\mathbb{Z}[X, Y]$?

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• F is homogeneous and $Y \nmid F$, so it is enough to consider

$$G(X) := F(X,1) = 3X^3 + 3X^2 + X^4 - 6$$

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► Eisenstein's criterion → G(X) is irreducible → F(X, Y) does not factor!

Let
$$p(z) = a_d z^d + ... + a_0$$
 be a polynomial, and let $p^{(n)}(z) = \underbrace{p \circ p \circ \cdots \circ p}_{n \text{ times}}(z)$. For certain k , compute

$$\lim_{n\to\infty}\frac{1}{d^n}\log\max(1,|p^{(n)}(k)|)\;.$$

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•
$$p(z) = z^d + t$$
 and $k = 0$ is tractable for a computer

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Count the number of \mathbb{Z}^2 lattice points in the region bounded by y = 2/3x + 4, y = 3/2x - 4, y = -2x + 2, y = -1/3x - 5.

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$$p(z) = z^2 + c$$
 for $c \in \mathbb{C}$, and define
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Question: How does g(c) vary in c?

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Consider the following two subgroups of S_7 :

 $G_1 = \{ \mathrm{id}, (34)(65), (47)(16), (37)(15), (156)(473), (165)(374) \},\$

 $G_2 = {id, (12)(56), (25)(16), (26)(15)}.$

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$$\begin{split} G_1 &= \{ \mathrm{id}, (34)(65), (47)(16), (37)(15), (156)(473), (165)(374) \}, \\ G_2 &= \{ \mathrm{id}, (12)(56), (25)(16), (26)(15) \}. \end{split}$$

Question: How many elements does the subgroup of S_7 generated by G_1 and G_2 have?

Consider the elliptic curve given in Weierstrass form by

$$E: y^2 = x^3 + 8 .$$

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Question:

- ▶ Over Q, does the point *P*(1,3) have infinite or finite order?
- ► Over Q(√2), does the point Q(0, 2√2) have infinite or finite order on E?

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In \mathbb{R}^5 , let *C* be the cone with vertex at the origin and whose extremal rays are generated by

$$(10, 1, 3, -1, 4), (-3, 0.5, 5, 7, 2), (1, 6, 3, 7, 11),$$

 $(9, 1, 1, 3, 5), (5, 0, 1, 4, -1), (6, 1, 2, 7, 3).$

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Question: Is the vector (1, 1, 2, 2, 3) in C or not?

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- Sage;
- ► GAP;
- Macaulay2;
- LRS, porta;
- Polymake;
- Python;
- ► C, C++.

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- Mathematica;
- Matlab;
- Maple;
- Magma;
- Geometer's Sketchpad;
- vlab.uga.edu

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- Your office computer (should have all licensed programs pre-installed)
- 6th floor computers faster processors, more memory; requires access from help@math.uga.edu
- UGA cluster (requires remote access)

Mathematica:

http://reference.wolfram.com/language/

LRS:

http://cgm.cs.mcgill.ca/~avis/C/lrslib/USERGUIDE.html

- Sage: http://doc.sagemath.org/html/en/
- GAP: http://www.gap-system.org/Doc/doc.html
- Polymake: https://polymake.org/doku.php/documentation
- Macaulay2: http://www.math.uiuc.edu/Macaulay2/
- Matlab: http://www.mathworks.com/help/matlab/
- Maple: http://www.maplesoft.com/documentation_center/
- Porta:

http://comopt.ifi.uni-heidelberg.de/software/PORTA/